Chapter 3
Transport Layer

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Computer Networking: A Top Down Approach
5th edition.
Jim Kurose, Keith Ross
Addison-Wesley, April 2009.
Chapter 3 outline

3.1 Transport-layer services
3.2 Multiplexing and demultiplexing
3.3 Connectionless transport: UDP
3.4 Principles of reliable data transfer

3.5 Connection-oriented transport: TCP
   - segment structure
   - reliable data transfer
   - flow control
   - connection management

3.6 Principles of congestion control
3.7 TCP congestion control
TCP reliable data transfer

- TCP creates rdt service on top of IP’s unreliable service
- pipelined segments
- cumulative acks
- TCP uses single retransmission timer

- retransmissions are triggered by:
  - timeout events
  - duplicate acks
- initially consider simplified TCP sender:
  - ignore duplicate acks
  - ignore flow control, congestion control
TCP sender events:

**data rcvd from app:**
- Create segment with seq #
- seq # is byte-stream number of first data byte in segment
- start timer if not already running (think of timer as for oldest unacked segment)
- expiration interval: TimeOutInterval

**timeout:**
- retransmit segment that caused timeout
- restart timer

**Ack rcvd:**
- If acknowledges previously unacked segments
  - update what is known to be acked
  - start timer if there are outstanding segments
NextSeqNum = InitialSeqNum
SendBase = InitialSeqNum

loop (forever) {
  switch(event)

  event: data received from application above
  create TCP segment with sequence number NextSeqNum
  if (timer currently not running)
    start timer
  pass segment to IP
  NextSeqNum = NextSeqNum + length(data)

  event: timer timeout
  retransmit not-yet-acknowledged segment with
  smallest sequence number
  start timer

  event: ACK received, with ACK field value of y
  if (y > SendBase) {
    SendBase = y
    if (there are currently not-yet-acknowledged segments)
      start timer
  }

} /* end of loop forever */

Comment:
  • SendBase-1: last cumulatively acked byte

Example:
  • SendBase-1 = 71; y= 73, so the rcvr wants 73+ ;
    y > SendBase, so that new data is acked
TCP: retransmission scenarios

Host A
Seq=100, 20 bytes data
ACK=100

Host B
Seq=92, 8 bytes data
Seq=92, 8 bytes data
Seq=92 timeout

Host A
Seq=92, 8 bytes data
ACK=100

Host B
Seq=92 timeout

SendBase = 100

SendBase = 100

SendBase = 120

SendBase = 120

lost ACK scenario

lost ACK scenario

premature timeout

X
loss

timeout

time

SendBase = 100

SendBase = 100

SendBase = 120

SendBase = 120

Seq=92, 8 bytes data
Seq=92 timeout

Seq=100, 20 bytes data
ACK=100

Seq=92, 8 bytes data

ACK=120

ACK=120

Seq=92 timeout

premature timeout
TCP retransmission scenarios (more)
## TCP ACK generation [RFC 1122, RFC 2581]

<table>
<thead>
<tr>
<th>Event at Receiver</th>
<th>TCP Receiver action</th>
</tr>
</thead>
<tbody>
<tr>
<td>Arrival of in-order segment with expected seq #. All data up to expected seq # already ACKed</td>
<td>Delayed ACK. Wait up to 500ms for next segment. If no next segment, send ACK</td>
</tr>
<tr>
<td>Arrival of in-order segment with expected seq #. One other segment has ACK pending</td>
<td>Immediately send single cumulative ACK, ACKing both in-order segments</td>
</tr>
<tr>
<td>Arrival of out-of-order segment higher-than-expect seq. # . Gap detected</td>
<td>Immediately send <em>duplicate ACK</em>, indicating seq. # of next expected byte</td>
</tr>
<tr>
<td>Arrival of segment that partially or completely fills gap</td>
<td>Immediate send ACK, provided that segment starts at lower end of gap</td>
</tr>
</tbody>
</table>
Fast Retransmit

- time-out period often relatively long:
  - long delay before resenting lost packet
- detect lost segments via duplicate ACKs.
  - sender often sends many segments back-to-back
  - if segment is lost, there will likely be many duplicate ACKs.
- if sender receives 3 ACKs for the same data, it supposes that segment after ACKed data was lost:
  - fast retransmit: resend segment before timer expires
Figure 3.37 Resending a segment after triple duplicate ACK
Fast retransmit algorithm:

event: ACK received, with ACK field value of \( y \)
  if \( y > \text{SendBase} \) {
    \text{SendBase} = y
    if (there are currently not-yet-acknowledged segments)
      start timer
  }
else {
  increment count of dup ACKs received for \( y \)
  if (count of dup ACKs received for \( y = 3 \)) {
    resend segment with sequence number \( y \)
  }
}

a duplicate ACK for already ACKed segment

fast retransmit
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TCP Flow Control

- receive side of TCP connection has a receive buffer:
  - speed-matching service: matching the send rate to the receiving app’s drain rate
  - flow control: sender won’t overflow receiver’s buffer by transmitting too much, too fast

- app process may be slow at reading from buffer
TCP Flow control: how it works

- rcvr advertises spare room by including value of RcvWindow in segments
- sender limits unACKed data to RcvWindow
  - guarantees receive buffer doesn’t overflow

(suppose TCP receiver discards out-of-order segments)

- spare room in buffer
  - \( \text{RcvWindow} = \text{RcvBuffer} - [\text{LastByteRcvd} - \text{LastByteRead}] \)
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TCP Connection Management

**Recall:** TCP sender, receiver establish “connection” before exchanging data segments

- initialize TCP variables:
  - seq. #s
  - buffers, flow control info (e.g. RcvWindow)

- **client:** connection initiator
  ```java
  Socket clientSocket = new Socket("hostname","port number");
  ```

- **server:** contacted by client
  ```java
  Socket connectionSocket = welcomeSocket.accept();
  ```

**Three way handshake:**

**Step 1:** client host sends TCP SYN segment to server
- specifies initial seq #
- no data

**Step 2:** server host receives SYN, replies with SYNACK segment
- server allocates buffers
- specifies server initial seq. #

**Step 3:** client receives SYNACK, replies with ACK segment, which may contain data
TCP Connection Management (cont.)

Closing a connection:

client closes socket:
   clientSocket.close();

**Step 1:** client end system sends TCP FIN control segment to server

**Step 2:** server receives FIN, replies with ACK. Closes connection, sends FIN.
TCP Connection Management (cont.)

**Step 3:** client receives FIN, replies with ACK.
- Enters “timed wait” - will respond with ACK to received FINs

**Step 4:** server, receives ACK. Connection closed.

**Note:** with small modification, can handle simultaneous FINs.
TCP Connection Management (cont)

TCP client lifecycle

TCP server lifecycle